

Exam

Exercise 1

Let $b \in \mathbb{R}^m$, $A \in \mathbb{R}^{m \times n}$. We want the solution of $\min \|x\|_\infty$ (P)
s.t. $Ax = b$

$$(P) \Leftrightarrow \min_{t, x} t \quad \Leftrightarrow \min_{t, x} t$$

s.t. $\|x\|_\infty \leq t$ s.t. $|x_1| \leq t, \dots, |x_n| \leq t$
 $Ax = b$ $Ax = b$

$$\Leftrightarrow \min_{t, x} t \quad \Leftrightarrow \min_{t, x} t$$

s.t. $-t \leq x_1 \leq t, \dots, -t \leq x_n \leq t$ s.t. $x \leq t \mathbb{1}_n$
 $Ax = b$ $x \geq -t \mathbb{1}_n$
 $Ax = b$

Therefore (P) can be rewritten as a LP : $\min_{t, x} t$ (LP)
s.t. $x \leq t \mathbb{1}_n$
 $-x \leq -t \mathbb{1}_n$
 $Ax = b$

The associated Lagrangian is:

$$L(t, x, \lambda_1, \lambda_2, \nu) = t + \lambda_1^T (x - t \mathbb{1}_n) + \lambda_2^T (-x - t \mathbb{1}_n) + \nu^T (Ax - b)$$
$$= t(1 - \lambda_1^T \mathbb{1}_n - \lambda_2^T \mathbb{1}_n) + (\lambda_1^T - \lambda_2^T + \nu^T A)x - \nu^T b$$

The dual function is $g(\lambda_1, \lambda_2, \nu) = \begin{cases} -\nu^T b & \text{if } \begin{cases} \lambda_1 + \lambda_2 + \nu^T A = 0 \\ \lambda_1 - \lambda_2 + \nu^T \mathbb{1}_n = 1 \end{cases} \\ -\infty & \text{otherwise} \end{cases}$

And the dual problem is : $\max g(\lambda_1, \lambda_2, \nu)$
 s.t. $\lambda_1 \geq 0$
 $\lambda_2 \geq 0$

which is equivalent to $\max -\nu^T b$
 s.t. $\lambda_1 \geq 0, \lambda_2 \geq 0$
 $(\lambda_1 + \lambda_2)^T \mathbf{1}_n = 1$
 $\lambda_1 - \lambda_2 + A^T \nu = 0$

We can finally rewrite the dual problem as:

$$\begin{aligned} \max & -\nu^T b \\ \text{s.t.} & \lambda \geq 0 \\ & \lambda + A^T \nu \geq 0 \\ & \underline{(\lambda + A^T \nu)^T \mathbf{1}_n = 1} \end{aligned}$$

Exercise 2

Let $u, v \in \mathbb{R}_+^m$. $D_{\text{KL}}(u, v) = \sum_{i=1}^m u_i \log\left(\frac{u_i}{v_i}\right) - u_i + v_i$

If one notes $f(u) = \sum_{i=1}^m u_i \log(u_i)$ the negative entropy of u , one has:
 $\nabla f(u) = \sum_{i=1}^m (\log(u_i) + 1) e_i$ on \mathbb{R}_+^m where $(e_i)_{i \in \{1, \dots, m\}}$ is the canonical basis of \mathbb{R}^n .

$$\begin{aligned} \text{Therefore } f(u) - f(v) - \nabla f(v)^T (u - v) &= \\ &= \sum u_i \log(u_i) - \sum v_i \log(v_i) - \sum (\log(v_i) + 1)(u_i - v_i) \\ &= \sum u_i (\log(u_i) - \log(v_i)) - u_i + v_i + 0 \\ &= D_{\text{KL}}(u, v) \end{aligned}$$

Moreover one has that $h: x \mapsto x \log(x)$ is \checkmark strictly convex for $x \in \mathbb{R}_+$

(because $h''(x) = 1/x > 0$) so f is strictly convex (as a sum of strictly convex functions)

Moreover, f is differentiable and therefore the first order condition of strict convexity gives that $f(u) > f(v) + \nabla f(v)^T(u-v)$

$$\Rightarrow \underline{D_u f(u,v) = f(u) - f(v) - \nabla f(v)^T(u-v) > 0}$$

Finally, if $u=v \in \mathbb{R}_{++}^n$, then $\underline{D_u f(u,v) = 0}$

Hence: $D_u f(u,v) \geq 0 \quad \forall u,v \in \mathbb{R}_{++}^n$

$$\underline{D_u f(u,v) = 0 \Leftrightarrow u=v}$$

Exercise 3

Let $A \in S_n$, $b \in \mathbb{R}^n$, $c \in \mathbb{R}^n$, $D \in \mathbb{R}^{m \times m}$, $g \in \mathbb{R}^m$

$A \succeq 0$ convex

and (P) : $\min c^T x$

$$\text{s.t. } x^T (A - bb^T) x \leq 0$$

$$b^T x \geq 0$$

$$Dx = g$$

Since $A \in S_n$ and $A \succeq 0$, there exists $V \in S_n$ such that $A = V^T V$.

Indeed, by the spectral theorem, there exists O orthogonal and Σ diagonal with positive values such that $A = O^T \Sigma O$.

The matrix $V = O^T \Sigma^{1/2} O$ works

Then one can rewrite the condition $x^T(A - bb^T)x \leq 0$
 as $x^T V^T V x - x^T b b^T x \leq 0 \Leftrightarrow \|Vx\|_2^2 \leq (b^T x)^2$

Since the condition $b^T x > 0$ holds, the above inequality
 is equivalent to $\|Vx\|_2 \leq b^T x$

Therefore, we can rewrite (P) as a SOCP:

$$\begin{aligned} \min \quad & c^T x && \text{(SOCP) which is a convex problem} \\ \text{s.t.} \quad & \|Vx\|_2 \leq b^T x \\ & Dx = g \end{aligned}$$

Now introduce a new variable $y = Vx$

$$\begin{aligned} \text{(SOCP)} \Rightarrow \min \quad & c^T x \\ \text{s.t.} \quad & \|y\|_2 \leq b^T x \\ & Vx = y \\ & Dx = g \end{aligned}$$

The Lagrangian is (for $\lambda \in \mathbb{R}$, $v_1 \in \mathbb{R}^m$, $v_2 \in \mathbb{R}^m$)

$$\begin{aligned} L(x, y, \lambda, v_1, v_2) &= c^T x + \lambda (\|y\|_2 - b^T x) + v_1^T (Vx - y) + v_2^T (Dx - g) \\ &= (c - b\lambda + V^T v_1 + D^T v_2)^T x + \lambda \|y\|_2 - v_1^T y - v_2^T g \end{aligned}$$

Now recall that the conjugate of $x \mapsto \|x\|_2$ is $y \mapsto \begin{cases} 0 & \text{if } \|y\|_2 \leq 1 \\ \infty & \text{otherwise} \end{cases}$
 and that the ℓ_2 norm is self dual.

$$\begin{aligned} \text{Then } \inf_y (\lambda \|y\|_2 - v_1^T y) &= -\lambda \sup \left(\frac{1}{\lambda} v_1^T y - \|y\|_2 \right) \quad (\lambda \neq 0) \\ &= \begin{cases} 0 & \text{if } \|v_1\|_2 \leq |\lambda| \\ -\infty & \text{otherwise} \end{cases} \end{aligned}$$

The obtained formula also works for the case $\lambda = 0$.

Therefore the dual function is:

$$g(\lambda, v_1, v_2) = \begin{cases} -v_2^T g & \text{if } \begin{cases} c - b\lambda + v_1^T v_1 + D^T v_2 = 0 \\ \|v_1\|_2 \leq |\lambda| \end{cases} \\ -\infty & \text{otherwise} \end{cases}$$

The dual problem can be written as:

$$\begin{aligned} \max g(\lambda, v_1, v_2) &\Leftrightarrow \max -v_2^T g \\ \text{s.t. } \lambda > 0 &\quad \text{s.t. } \begin{cases} c - b\lambda + v_1^T v_1 + D^T v_2 = 0 \\ \|v_1\|_2 \leq |\lambda| \\ \lambda \geq 0 \end{cases} \end{aligned}$$

$$\Leftrightarrow \begin{aligned} \max & -v_2^T g \\ \text{s.t. } & \begin{cases} c - b\lambda + v_1^T v_1 + D^T v_2 = 0 \\ \|v_1\|_2 \leq \lambda \end{cases} \end{aligned}$$

Exercise 4

Let (P): $\min -\sum_{i=1}^m \log(b_i - a_i^T x)$ over $\{x \mid a_i^T x < b_i, i=1, \dots, m\}$

Introduce a new variable $y = b - Ax$ where the i th row of A is a_i^T , so that $y_i = b_i - a_i^T x$.

The new problem is $\min \sum_{i=1}^m \log(y_i)$
 s.t. $y = b - Ax$

$$x \in \mathcal{D} \Leftrightarrow y \in \mathbb{R}_{++}^m$$

The associated Lagrangian is (for $v \in \mathbb{R}^m$, and $x \in \mathcal{D}, y \in \mathbb{R}_{++}^m$)

$$\begin{aligned} L(x, y, v) &= -\sum_{i=1}^m \log(y_i) + v^T (y - b + Ax) \\ &= -\sum_{i=1}^m \log(y_i) + v^T y + v^T A x - v^T b \end{aligned}$$

Let $f: y \in \mathbb{R}_{++}^m \mapsto -\sum_{i=1}^m \log y_i + v^T y$. We suppose $v_i > 0 \quad \forall i=1, \dots, m$

On its domain, f is differentiable and one has:

$(\nabla f(y))_i = v_i - 1/y_i$, which is differentiable.

The hessian of f is $Hf(y) = \begin{pmatrix} -1/y_1^2 & & \\ & \ddots & \\ & & -1/y_m^2 \end{pmatrix} \prec 0$

Therefore f is convex and is such that $\nabla f(y_0) = 0$ for

$$y_0 = (1/v_1, \dots, 1/v_m)^T$$

So in the case $v > 0$, f reaches a minimum in y_0

$$\text{that is } \min_{y \in \mathbb{R}_{++}^m} f(y) = m + \sum_{i=1}^m \log(v_i)$$

If $v \not> 0$: $\exists i=1, \dots, m$ s.t. $v_i \leq 0$

then let $y_t = (1, \dots, 1, t, 1, \dots, 1)$ ($t > 0$)
 \uparrow
 i -th entry

$$f(y_t) = -\log(t) + v_i t + c$$

$\rightarrow -\infty$ as $t \rightarrow +\infty$ so f is unbounded below.

Therefore, the dual function is:

$$g(v) = \begin{cases} m + \sum_{i=1}^m \log(v_i) - v^T b & \text{if } \begin{cases} v^T A = 0 \\ v > 0 \end{cases} \\ -\infty & \text{otherwise} \end{cases}$$

And the dual problem reads:

$$\begin{aligned} \max \quad & m + \sum_{i=1}^m \log(v_i) - v^T b \\ \text{s.t.} \quad & v^T A = 0 \\ & v > 0 \end{aligned}$$

Exercise 5

$$(1) : \min f(x) \quad , \quad \phi(x) = f(x) + \alpha \|Ax - b\|_2^2$$

s.t. $Ax = b$

($f: \mathbb{R}^n \rightarrow \mathbb{R}$ convex and differentiable, $A \in \mathbb{R}^{m \times n}$, $\text{rank}(A) = m$)

We suppose \tilde{x} minimizes ϕ . Since f is differentiable, so is ϕ .

$$\text{Therefore } \nabla \phi(\tilde{x}) = 0 \Rightarrow \nabla f(\tilde{x}) + 2\alpha A^T(A\tilde{x} - b) = 0$$
$$\Rightarrow \nabla f(\tilde{x}) + A^T \underbrace{(2\alpha(A\tilde{x} - b))}_{=\tilde{v}} = 0$$

The Lagrangian of (1) is $L(x, v) = f(x) + v^T(Ax - b)$
which is convex and differentiable.

$$\nabla_x L(x, v) = \nabla f(x) + Av$$

Hence, for $\tilde{v} = 2\alpha(A\tilde{x} - b)$ is dual feasible and

$$g(\tilde{v}) = f(\tilde{x}) + \tilde{v}^T(A\tilde{x} - b)$$
$$= \underbrace{f(\tilde{x}) + 2\alpha \|A\tilde{x} - b\|_2^2}_{\phi(\tilde{x})} \leq f(x) \quad \forall x, Ax = b$$

$g(\tilde{v})$ is a lower bound on the optimal value of (1)

Exercise 6

See Jupyter Notebook